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#Input:
str1 = input("Enter string: ")
substr1 = input("Enter substring: ")
occur1 = input("Enter character whose occurrence is to be calculated: ")
occur2 = str1.split()
index1 = []
for i in occur2:
    for j in occur2:
        if(i == j and j not in index1):
            index1.append(i)
        else:
            continue

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def longestWord():
    globalMax = 0
    currentMax = 0
    list1 = []
    for i in str1:
        if (i != " "):
            currentMax = currentMax + 1
            list1.append(i)
        else:
            if (currentMax > globalMax):
                globalMax = currentMax
                temp = list1
            currentMax = 0
            list1 = []
    if (i == str1[len(str1) - 1]):
        if (currentMax > globalMax):
            globalMax = currentMax
            temp = list1
    print("Longest word is: ", "".join(temp))
    print("Length of largest word is: ", globalMax)

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def palindrome():
    rev1 = str1
    if (str1[::-1] == rev1):
        print("String is palindrome")
    else:
        print("String is not palindrome")

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def charfrequency():
    temp = 0
    for i in str1:
        max1 = 0
        for j in range(0, len(str1)):
            if (str1[j] != " " and str1[j] == occur1):
                max1 = max1 + 1
            else:
                continue

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    if (max1 > temp):
        temp = max1
    else:
        continue
print("Number of times the character '" + occur1, "' occurs: ", temp)

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def indexSubstring():
    for i in range(len(str1) - len(substr1) + 1):
        if str1[i:i + len(substr1)] == substr1:
            print("Substring present in string at index: ", i)
            break
    else:
        print("Substring not present")

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def occurWord():
    for i in index1:
        max1 = 0
        for j in occur2:
            if(i == j):
                max1 = max1 + 1
            else:
                continue
        print("The number of times '" + i, "' is repeated is: ", max1)

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flag = 1
while flag == 1:
    print("1. Display word with longest word: \n")
    print("2. The occurrence of '" + occur1, "' in the string is: \n")
    print("3. Check whether the string is palindrome or not: \n")
    print("4. Display index of first appearance of the substring: \n")
    print("5. Occurrence of each word in the given string: \n")
    print("6. Exit")
    ch = int(input("Enter your choice: "))
    if(ch == 1):
        longestWord()
    elif(ch == 2):
        charfrequency()
    elif(ch == 3):
        palindrome()
    elif(ch == 4):
        indexSubstring()
    elif(ch == 5):
        occurWord()
    elif(ch > 6 or ch < 1):
        print("Enter valid choice")
    elif(ch == 6):
        flag = 0

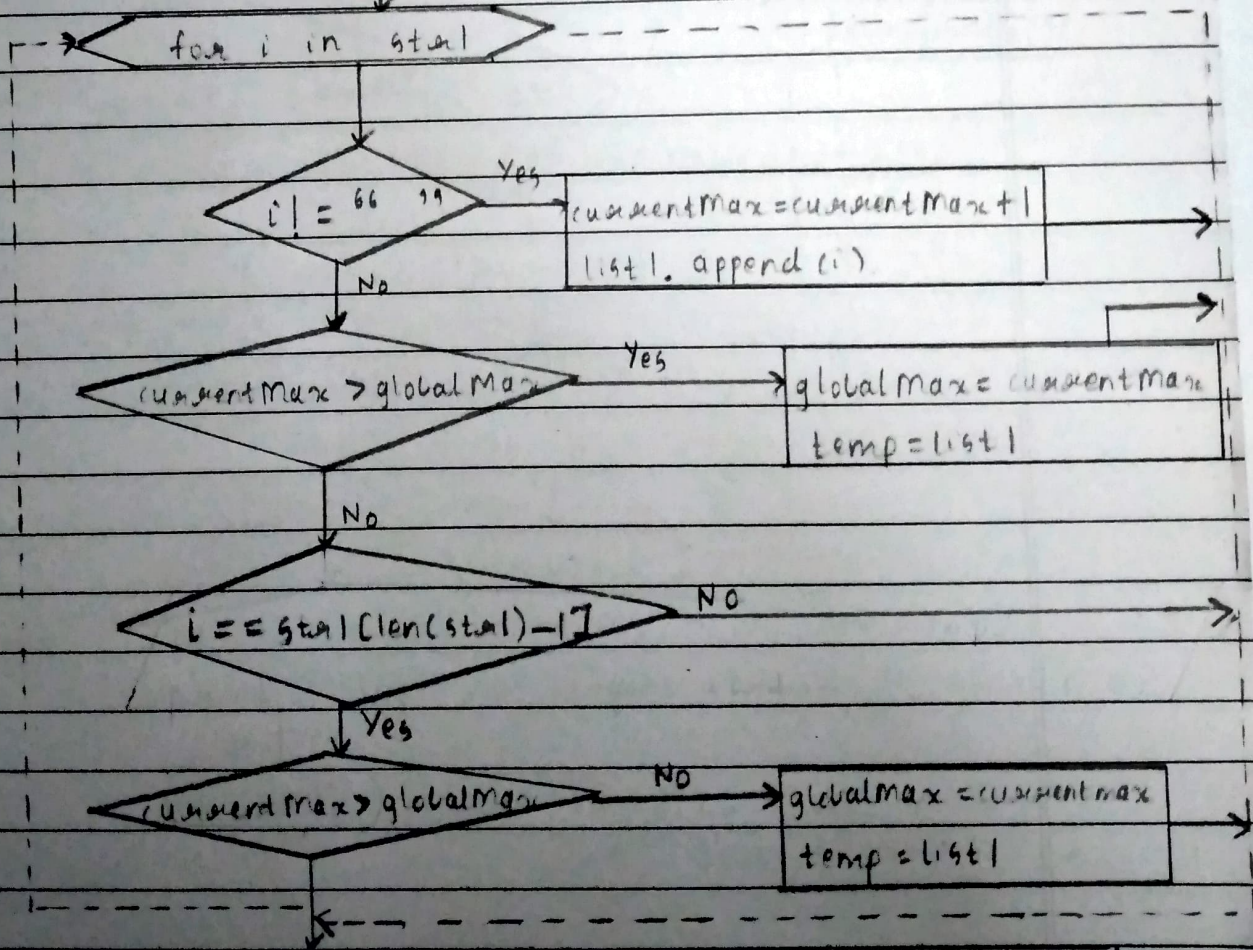
```


→ Flowchart for def longestWord():

~~def longestWord~~

def longestWord()

globalMax = 0
currentMax = 0
list1 = []

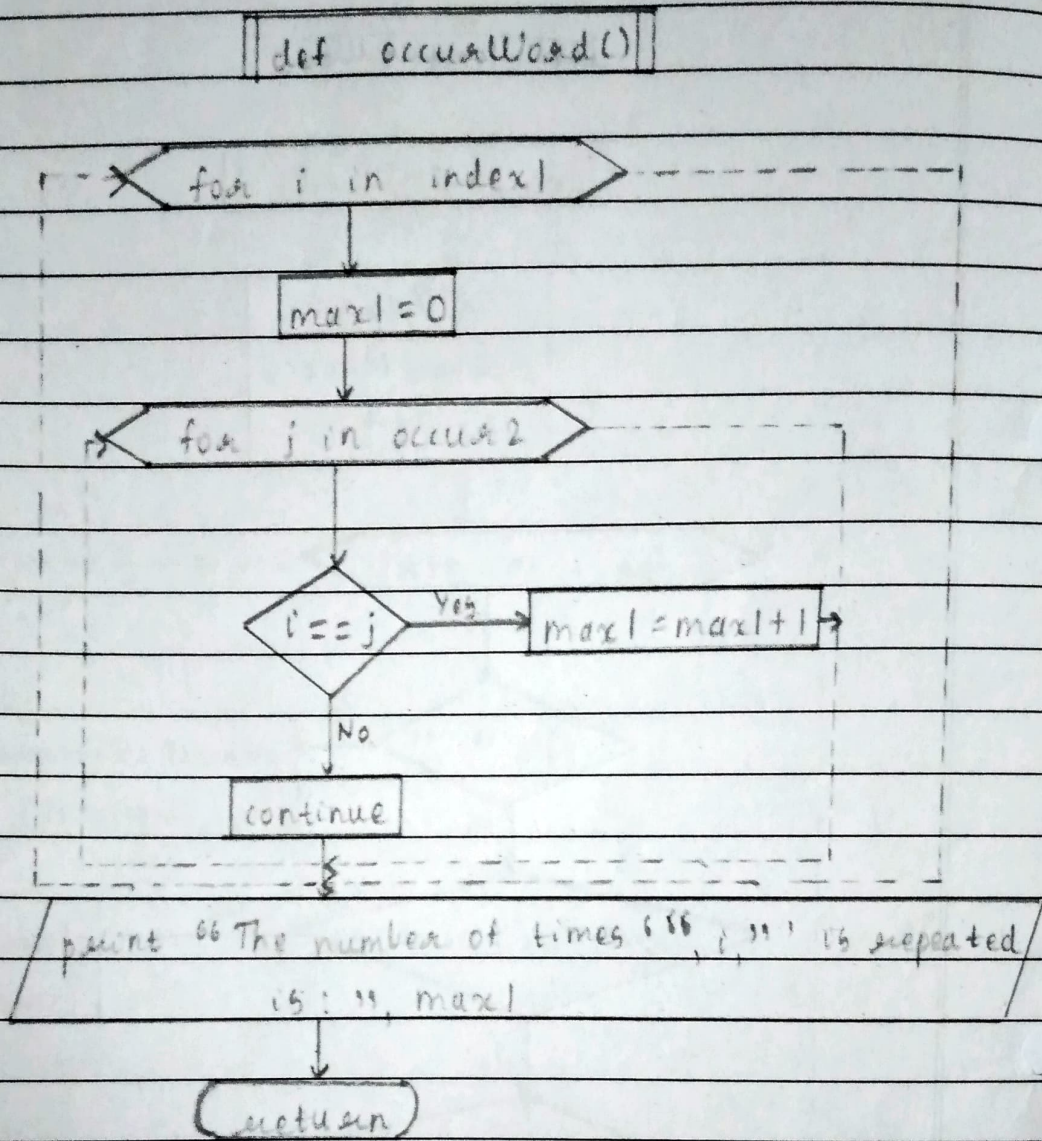


print "longest word is: ", join(temp)

print "Length of longest word is: ", globalMax

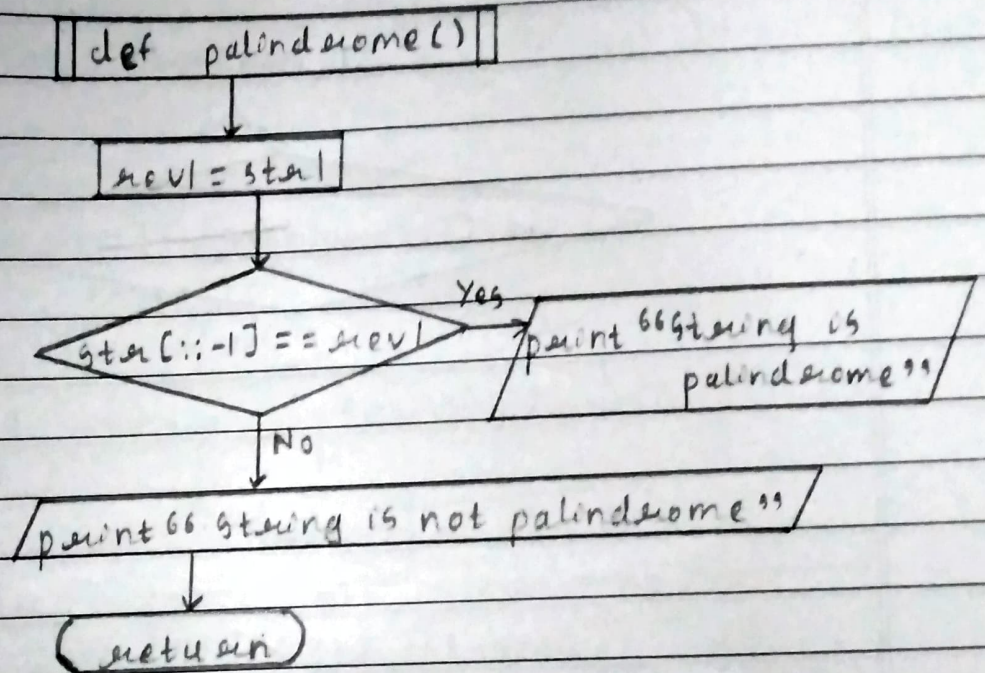
return

→ Flowchart for def occurWord()

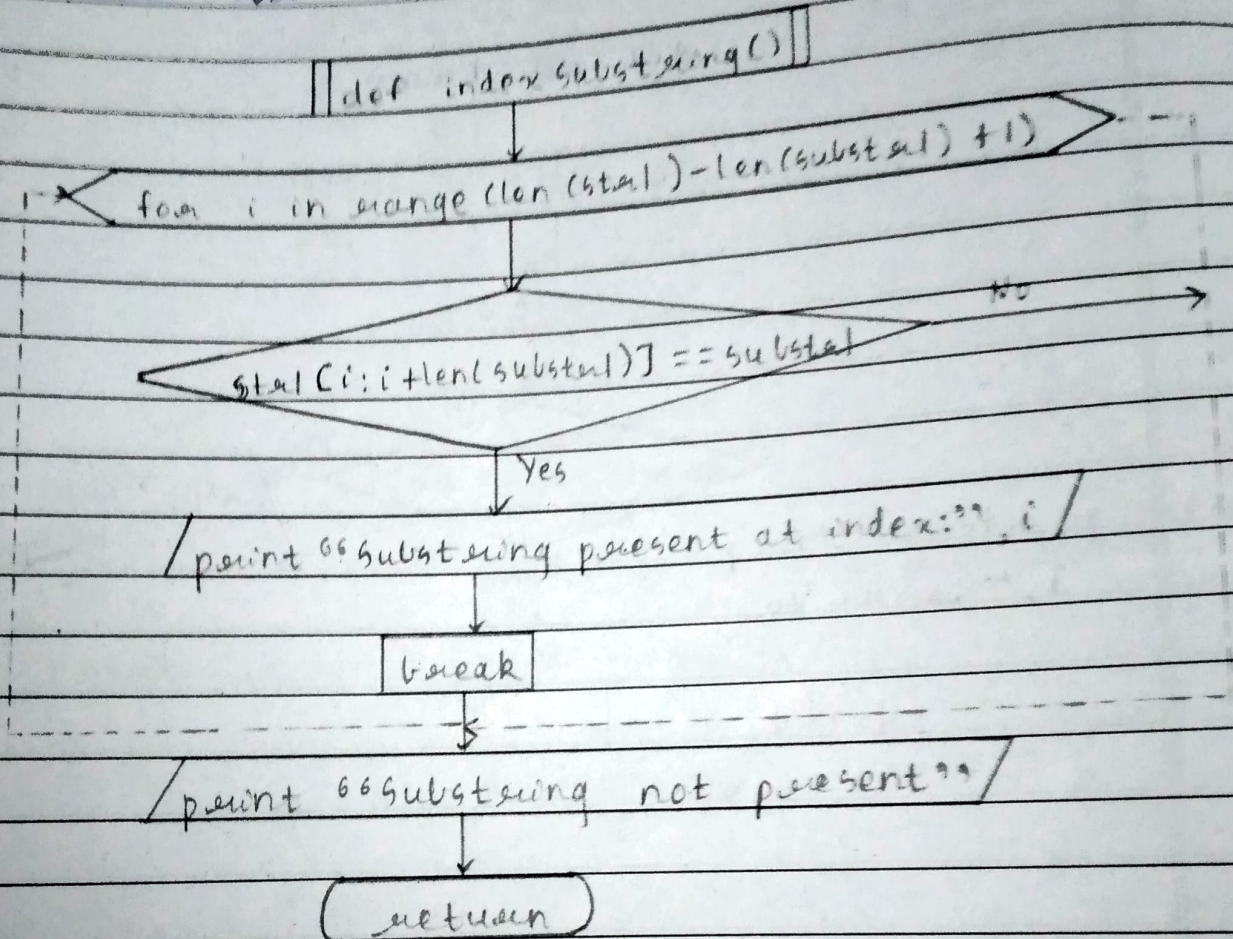


→ Flowchart for def palindrome()

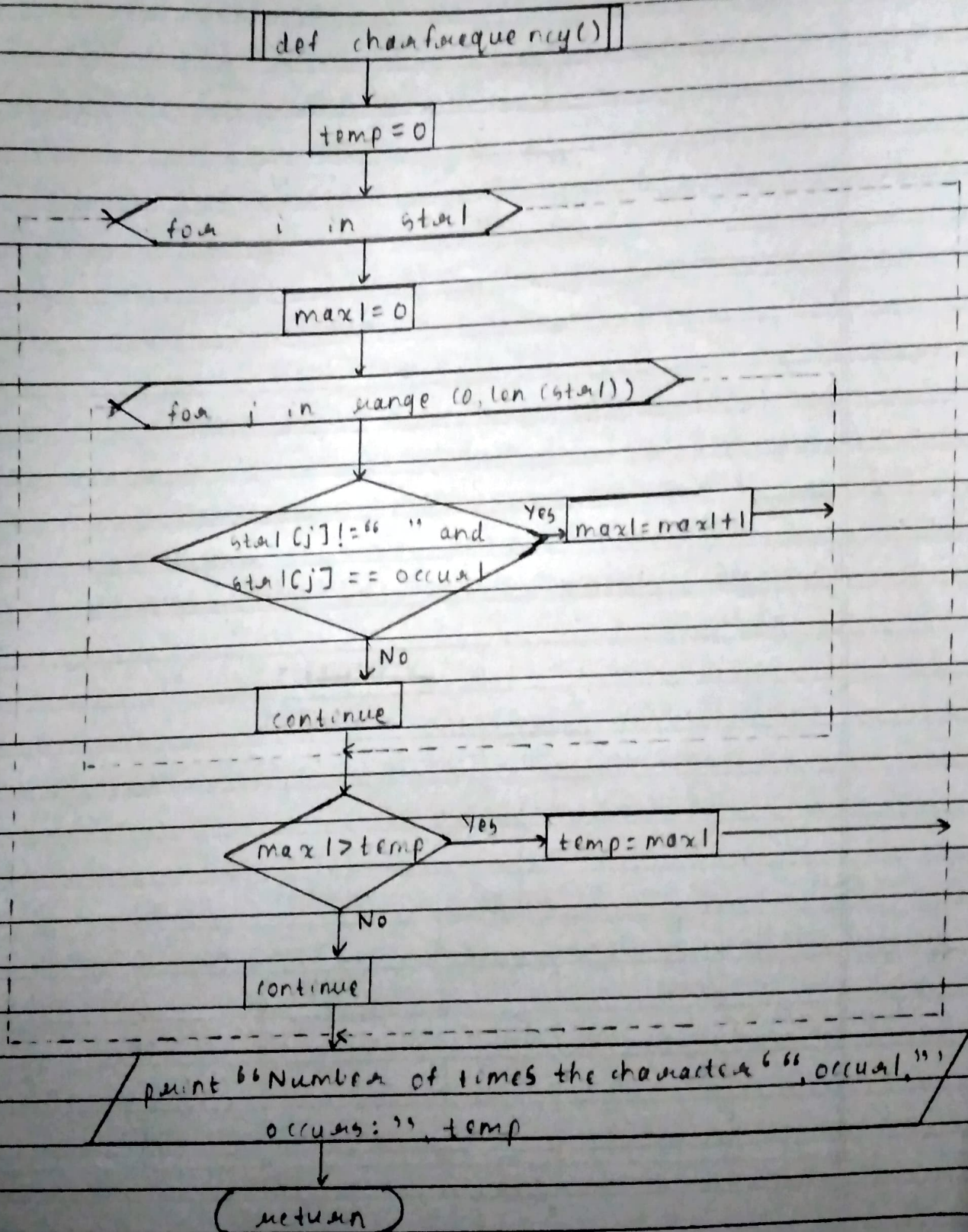
~~def palindrome~~



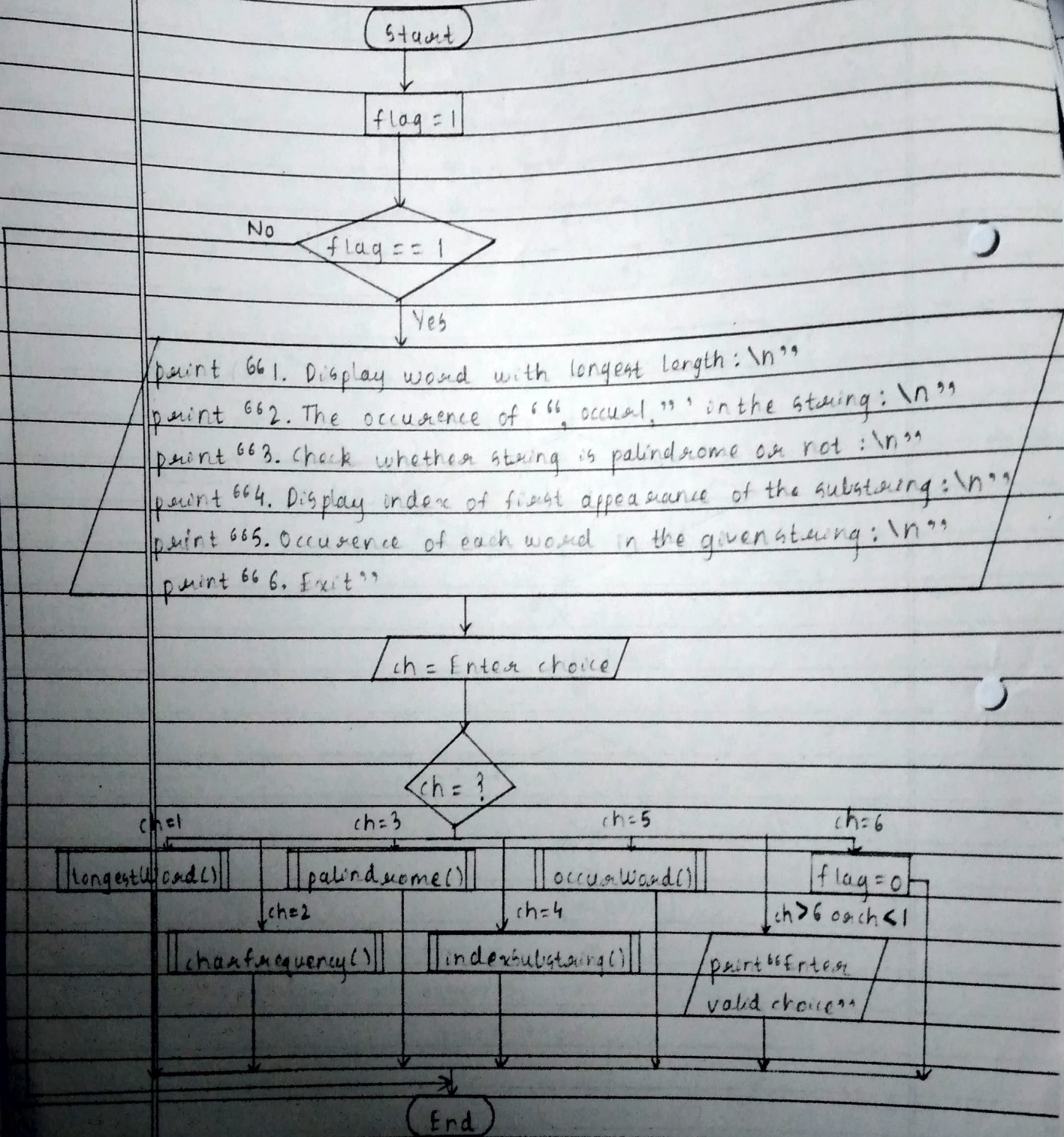
Flowchart for def indexOfSubstring()



→ Flowchart for def charfrequency()



→ Flowchart for menu



→ Pseudocode for def longestWord()

1. Initialize globalMax = 0

Initialize currentMax = 0

list1 = []

2. for i in str1 do

begin

if (ci != " ") then

currentMax = currentMax + 1

list1.append(ci)

else

if (currentMax > globalMax) then

globalMax = currentMax

store temp = list1

initialize currentMax = 0

list1 = []

if (ci == str1[len(str1) - 1]) then

if (currentMax > globalMax) then

globalMax = currentMax

store temp = list1

end

3. Display longest word is, temp

Display length of largest word is, globalMax

4. return

→ Pseudocode for def charFrequency()

1. Initialize temp=0

2. for i in str1 do
begin

 initialize max1=0

 for j in range (0, len (str1)) do

 begin

 if (str1[j] != " " and str1[j] == occur1) then

 max1 = max1 + 1

 else

 continue

 end

 if (max1 > temp) then

 store temp = max1

 else

 continue

end

3. Display Number of times the character, occur1,
 occur2, temp

→ Pseudocode for def palindrome

1. store rev1 = str1

2. if (str1[::-1] == rev1) then

 display string is palindrome

else

 Display string is not palindrome

3. return

→ Pseudocode for def indexSubstring()

1. for i in range($\text{len}(\text{str}) - \text{len}(\text{substr}) + 1$) do
begin

if $\text{str}[i : i + \text{len}(\text{substr})] = \text{substr}$ then

Display substring is present in string at
index: i

break

else

Display substring is not present

end

2. return

→ Pseudocode for def occurWord()

1. for i in index1 do

begin

initialize $\text{max1} = 0$

for j in occur2 do

begin

if $(i == j)$ then

$\text{max1} = \text{max1} + 1$

else

continue

end

end

2. Display The number of times, i is repeated is: max1

3. return

→ Pseudocode
Flowchart for menu

1. Start

2. Initialize flag = 1

3. while flag = 1 do

begin

Display 1. Display word with longest length:

Display 2. The occurrence of, occurred, in the string is:

Display 3. Check whether the string is palindrome.

or not:

Display 4. Display index of first appearance of the substring:

Display 5. Occurrence of each word in the given string

Display 6. Exit

ch = Enter your choice

if (ch == 1) then

call function longestWord()

elif (ch == 2) then

call function charFrequency()

elif (ch == 3) then

call function palindrome()

elif (ch == 4) then

call function indexSubstring()

elif (ch == 5) then

call function occurWord()

elif (ch > 6 or ch < 1) then

Display Enter valid choice

elif (ch == 6) then

initialize flag = 0

end

4. Stop

Q1. What is frequency count? Why is frequency count important in the analysis of algorithm?

Ans. Efficiency of an algorithm can be measured by inserting a counter in the algorithm in order to count the number of times the basic operation is executed that counter is called frequency count (FC).

- Frequency count in algorithm gives you that how many times your program is run that is it is dependent on the loop used in the program.
- However, the loop defines the time complexity of your program.
- So, it means that the frequency count defines the time complexity of the program.
- Also, the frequency count gives us the order of growth of an algorithm.

Q2. Write an algorithm to compute the sum of the digits of the given number. Justify that your algorithm satisfies all the given characteristics of an algorithm.

Ans. → Algorithm:-

1. Read $n =$ Enter number
2. Initialize $total = 0$
3. while ($n \neq 0$) do
begin
 store $total = total + n \% 10$
 store $n = n // 10$
4. Display total
5. Stop



- Justification:
- 1) Input: The given algorithm has a well defined input 'n'.
 - 2) Output: The algorithm has a well defined output 'total'.
 - 3) Finiteness: The given ~~number~~ algorithm terminates when the input 'n' is given; the looping is done as per the condition and the output 'total' is displayed.
 - 4) Definiteness: There are no complex and ambiguous steps in the above algorithm.
 - 5) Effectiveness: The above algorithm has an input statement, output statement, assigning statement, conditional statement and can be easily converted into a program hence satisfying the effectiveness property of the algorithm.

Q3. ~~Q3.~~ Write an algorithm to perform sparse matrix addition and state its time complexity.

Ans. Algorithm:-

1. Initialize $i=1, j=1, k=1$
2. Assign $L1 = c[0][2]$
Assign $L2 = d[0][2]$
3. Let $e[0][0] = c[0][0]$
let $e[0][1] = c[0][1]$
4. while $i \leq L1$ and $j \leq L2$ go to next step, otherwise go to step 21
5. if $c[i][0] = d[j][0]$ go to next step, otherwise go to step 15



6. If $c[i][i] == d[j][i]$ go to next step, otherwise go to step 9

7. Assign $e[k][0] = c[i][0]$
 $e[k][1] = c[i][1]$
 $e[k][2] = c[i][2] + d[j][2]$

8. increment i, j, k

9. if $c[i][i] > d[j][i]$ go to next step, otherwise go to step 12

10. Assign $e[k][0] = d[j][0]$
 $e[k][1] = d[j][1]$
 $e[k][2] = d[j][2]$

11. increment j, k

12. if $c[i][i] < d[j][i]$ go to next step, otherwise go to step 4

13. Assign $e[k][0] = ~~d[j][0]~~ c[i][0]$
 $e[k][1] = c[i][1]$
 $e[k][2] = c[i][2]$

14. increment i, k

15. if $c[i][0] > d[j][0]$ go to next step, otherwise go to step 18

16. Assign $e[k][0] = d[j][0]$
 $e[k][1] = d[j][1]$
 $e[k][2] = d[j][2]$

17. increment j, k

18. if $c[i][0] < d[j][0]$ go to next step, otherwise go to step 4

19. Assign $e[k][0] = c[i][0]$
 $e[k][1] = c[i][1]$
 $e[k][2] = c[i][2]$

20. increment i, k

21. while $j < l_2$ and $i > l_1$ go to next step,
otherwise go to step 24
22. Assign $e[k][0] = d[i][0]$
 $e[k][1] = d[i][1]$
 $e[k][2] = d[i][2]$
23. increment j, k
24. while $j > l_2$ and $i < l_1$ go to next step,
otherwise go to step 26
25. Assign $e[k][0] = c[i][0]$
 $e[k][1] = c[i][1]$
 $e[k][2] = c[i][2]$
24. increment i, k
26. Assign $e[0][2] = k - 1$
27. Print array e
28. stop

→ Time complexity:

• If l_1 and l_2 are number of non-zero elements in first and second sparse matrix respectively, then time complexity is $O(l_1 + l_2)$.