

SPPU-TE-COMP-CONTENT – KSKA Git

Find the most cost-effective path to reach the final state from initial state using A* Algorithm

Initial state

2	8	3
1	6	4
7		5

g=0 (Depth)
h=4

g=1	2	8	3	2	8	3	g=1	2	8	3	g=1
h=5	1	6	4	1		4	h=3	1	6	4	h=4
f=6		7	5	7	6	5	f=4	7	5		f=5

2	8	3	g=2	2	3	g=2	2	8	3	g=2	
	1	4	h=3	1	8	4	h=3	1	4	h=4	
7	6	5	f=6	7	6	5	f=5	7	6	5	f=6

8	3	2	8	3	2	3	2	3				
2	1	4	7	1	4	1	8	4	1	2	4	f=7
7	6	5	6	5	7	6	5	7	6	5		

1	2	3
	8	4
7	6	5

1	2	3
7	8	4
	6	5

1	2	3
8		4
7	6	5

final state

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Explain Admissibility and Consistency of Heuristic function in A* Algorithm

1. Admissibility of a Heuristic function

A heuristic function $h(n)$ is admissible if it never overestimates the true cost from a given state to the goal state i.e. $h(n) \leq h^*(n)$

where $h(n)$ is the heuristic estimate from node n to the goal. $h^*(n)$ is the actual cost of the shortest path from n to the goal.

This ensures that the heuristic is optimistic and does not mislead the algorithm into thinking a path is cheaper than it really is.

2. Consistency of heuristic function

A heuristic is consistent (or monotonic) if, for every node n , and every successor n' of n , the estimated cost to the goal from n should satisfy the following condition:

$$h(n) \leq c(n, n') + h(n')$$

where

- $h(n)$ is the heuristic estimate of the cost from node n to the goal

- $c(n, n')$ is the actual cost of reaching the successor node n' from node n .

- $h(n')$ is the heuristic estimate of the cost from node n' to the goal.

- Consistency ensures that the cost function $f(n) = g(n) + h(n)$ is non-decreasing along any path. This condition guarantees that once a node is expanded in A*, it will not need to be re-expanded, thus ensuring correctness & efficiency of the search.