L1 Context Free Grammar & Context Free Language

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Context Free Grammar is defined by 4 tuples as $G = \{ V, \Sigma, S, P \}$ where

V = Set of Variables or Non-Terminal Symbols

Σ = Set of Terminal Symbols

S = Start Symbol

P = Production Rule

Context Free Grammar has Production Rule of the form

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<u>Example:</u> For generating a language that generates equal number of a's and b's in the form $a^n b^n$, the Context Free Grammar wil be defined as

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$\rightarrow \underline{a}^3\underline{b}^3 \Rightarrow a^n b^n$

Questions????